

Fiat categorification of IS_n and F_n^*

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CSA 2016 Lisbon, PORTUGAL

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Objects of study

$$n \in \{1, 2, 3, \dots\}$$

$$\mathbf{n} := \{1, 2, \dots, n\}$$

S_n — the **symmetric group** on \mathbf{n}

IS_n — the **symmetric inverse semigroup** on \mathbf{n} (bijections between subsets of \mathbf{n})

I_n^* — the **dual symmetric inverse semigroup** on \mathbf{n} (bijections between quotients of \mathbf{n})

$F_n^* := S_n E(I_n^*)$ — the **maximal factorizable submonoid** of I_n^*

2-categories

Definition. A **2-category** is a category enriched over the monoidal category **Cat** of small categories (in the latter the monoidal structure is induced by the cartesian product).

This means that a 2-category \mathcal{C} is given by the following data:

- ▶ **objects** of \mathcal{C} ;
- ▶ **small categories** $\mathcal{C}(i, j)$ of morphisms;
- ▶ **bifunctorial composition** $\mathcal{C}(j, k) \times \mathcal{C}(i, j) \rightarrow \mathcal{C}(i, k)$;
- ▶ **identity objects** $\mathbb{1}_i \in \mathcal{C}(i, i)$;

which are subject to the obvious set of (strict) axioms.

Terminology and the first example

Terminology:

- ▶ An **object** in $\mathcal{C}(i, j)$ is called a **1-morphism**
- ▶ A **morphism** in $\mathcal{C}(i, j)$ is called a **2-morphism**
- ▶ **Composition** in $\mathcal{C}(i, j)$ is called **vertical**, denoted \circ_1
- ▶ **Composition** in \mathcal{C} is called **horizontal**, denoted \circ_0

Principal example. The category **Cat** is a 2-category.

- ▶ Objects of **Cat** are **small categories**
- ▶ 1-morphisms in **Cat** are **functors**
- ▶ 2-morphisms in **Cat** are **natural transformations**
- ▶ Composition is the usual **composition**
- ▶ Identity 1-morphisms are **identity functors**

Example from semigroups

Example from ordered monoids. Let $\mathbf{S} := (S, \cdot, e, \leq)$ be an ordered monoid.

Define the 2-category $\mathcal{C}_{\mathbf{S}}$ as follows:

- ▶ $\mathcal{C}_{\mathbf{S}}$ has one (formal) object \mathbf{i} ;
- ▶ objects in $\mathcal{C}_{\mathbf{S}}(\mathbf{i}, \mathbf{i})$ are elements in S ;
- ▶ composition of 1-morphisms in $\mathcal{C}_{\mathbf{S}}$ is \cdot ;
- ▶ the identity 1-morphism in $\mathcal{C}_{\mathbf{S}}(\mathbf{i}, \mathbf{i})$ is e ;
- ▶ the set of 2-morphisms from s to t is empty if $s \not\leq t$ and consists of one element $m_{s,t}$ if $s \leq t$;
- ▶ vertical composition of 2-morphisms is the only possible map which exists due to transitivity of \leq ;
- ▶ horizontal composition of 2-morphisms is the only possible map which exists due to admissibility of \leq ;
- ▶ the identity 2-morphisms are $m_{s,s}$, $s \in S$.

Consequences

Consequence 1. For any **monoid** (S, \cdot, e) , the equality relation $=$ is an admissible order. This gives rise to the 2-category \mathcal{C}_S , where $\mathbf{S} := (S, \cdot, e, =)$.

Consequence 2. For any **inverse monoid** $(S, \cdot, e, ()^{-1})$, we have the **natural** partial order \prec , which is admissible. This also gives rise to the 2-category \mathcal{C}_S , where $\mathbf{S} = (S, \cdot, e, \prec)$.

Note. If S is inverse, then \mathcal{C}_S has, usually, **more** 2-morphisms than \mathcal{C}_S .

Note. Both constructions apply to S_n , IS_n and F_n^* .

Question. What are **disadvantages** of these constructions.

Adjoint functors

\mathcal{A}, \mathcal{C} — two categories

$F : \mathcal{A} \rightarrow \mathcal{C}, \quad G : \mathcal{C} \rightarrow \mathcal{A}$ — two functors

Definition. The pair (F, G) is an **adjoint pair of functors** provided that

there exist $\alpha : \text{Id}_{\mathcal{A}} \rightarrow GF$ and $\beta : FG \rightarrow \text{Id}_{\mathcal{C}}$

such that

$$(\beta \circ_0 F) \circ_1 (F \circ_0 \alpha) = \text{id}_F \quad \text{and} \quad (G \circ_0 \beta) \circ_1 (\alpha \circ_0 G) = \text{id}_G.$$

Note: In **Cat** this is defined purely in terms of 2-morphisms.

2-categories with weak involution

\mathcal{C} — 2-category

$\star : \mathcal{C} \rightarrow \mathcal{C}$ — weak **anti-autoequivalence** reversing the order of both 1-morphisms and 2-morphism

Weak: $(FG)^\star \cong G^\star F^\star$, not necessarily $(FG)^\star = G^\star F^\star$

Definition \mathcal{C} is **iat** provides that, for each 1-morphism F , there exist 2-morphisms making (F, F^\star) into an adjoint pair of 1-morphisms

iat: **i**nvolution, **a**djunction, **t**wo= 2-category

Our examples: \mathcal{S}_{S_n} is iat, while \mathcal{S}_{IS_n} and $\mathcal{S}_{F_n^\star}$ are not.

Why: Not enough 2-morphisms between 1-morphisms and the identity 1-morphism.

Main problem

Problem: Is it possible to “enlarge” \mathcal{S}_{IS_n} and $\mathcal{S}_{F_n^*}$ to something iat?

Answer: YES

Rest of the talk: Construction.

Rough idea

Step 1. Start with the **iat** 2-category \mathcal{S}_{S_n} .

Step 2. Enlarge \mathcal{S}_{S_n} (in different ways, depending on \mathcal{S}_{IS_n} or $\mathcal{S}_{F_n^*}$) by adding new 2-morphisms.

Step 3. Linearize (e.g. over \mathbb{Z} or \mathbb{C}).

Step 4. Split idempotents.

Enlarging \mathcal{S}_{S_n} in the case of IS_n

Define a 2-category \mathcal{A} as follows:

- ▶ it has one object \mathbf{i} ;
- ▶ its 1-morphisms are elements in S_n ;
- ▶ the composition of 1-morphisms is multiplication in S_n ;
- ▶ the identity 1-morphism is $\text{id}_n \in S_n$;
- ▶ for $\pi, \sigma \in S_n$, the set $\text{Hom}_{\mathcal{A}}(\pi, \sigma)$ is the set of all $\alpha \in \mathbf{B}_n$ (binary relations) such that $\alpha \subseteq \pi \cap \sigma$;
- ▶ for $\pi, \sigma, \tau \in S_n$, and also for $\alpha \in \text{Hom}_{\mathcal{A}}(\pi, \sigma)$ and $\beta \in \text{Hom}_{\mathcal{A}}(\sigma, \tau)$, we set $\beta \circ_1 \alpha := \beta \cap \alpha$;
- ▶ for $\pi \in S_n$, we define the identity element in $\text{Hom}_{\mathcal{A}}(\pi, \pi)$ to be π ;
- ▶ for $\pi, \sigma, \tau, \rho \in S_n$, and also for $\alpha \in \text{Hom}_{\mathcal{A}}(\pi, \sigma)$ and $\beta \in \text{Hom}_{\mathcal{A}}(\tau, \rho)$, we define $\beta \circ_0 \alpha := \beta\alpha$, the usual composition of binary relations.

What happened?

Theorem. The construct \mathcal{A} is a **2-category**.

Observation 1. We have $\text{Hom}_{\mathcal{A}}(\pi, \sigma) = \text{Hom}_{\mathcal{A}}(\sigma, \pi)$, for all π and σ , in particular, we often have **many morphisms to and from identity**.

Observation 2. We have $\text{End}_{\mathcal{A}}(\text{id}_{\mathbf{n}}) = E(IS_n)$, all binary relations which are subrelations of the identity. In particular, $\text{End}_{\mathcal{A}}(\text{id}_{\mathbf{n}})$ is commutative and has **many idempotents**.

Observation 3. \mathcal{S}_{S_n} is a subcategory of \mathcal{A} , in particular, \mathcal{A} is **iat**. Note that \mathcal{S}_{S_n} is not 2-full in \mathcal{A} .

What next?

Next step 1. Linearize 2-morphisms over \mathbb{C} (consider 2-morphisms as a basis for a \mathbb{C} -vector space and extend composition by bilinearity).

Next step 2. Take the additive closure on the level of 1-morphisms by adding formal direct sums.

Next step 3. Split idempotents on the level of 2-morphisms.

Outcome: a new 2-category, call it \mathcal{C} .

Properties: Finitarity: finitely many 1-morphisms up to iso and finite dimensional spaces of 2-morphisms. Inherited iat-ness.

Theorem. The 2-category \mathcal{C} is fiat.

Decategorification

Grothendieck decategorification. Consider the split Grothendieck ring $[\mathcal{C}(i, i)]_{\oplus}$.

As abelian group: Free abelian group on $[F]$, where F is a 1-morphism, modulo $[F] = [G] + [H]$ whenever $F \cong G \oplus H$.

Ring structure: Inherited from composition of 1-morphisms.

Complexify: $\mathbb{C} \otimes_{\mathbb{Z}} [\mathcal{C}(i, i)]_{\oplus}$.

Main Theorem. $\mathbb{C} \otimes_{\mathbb{Z}} [\mathcal{C}(i, i)]_{\oplus} \cong \mathbb{C}[IS_n]$, where the basis of indecomposable 1-morphisms corresponds to the **Möbius basis** of $\mathbb{C}[IS_n]$ (cf. B. Steinberg. Möbius functions and semigroup representation theory. J. Combin. Theory Ser. A **113** (2006), no. 5, 866–881.)

Enlarging \mathcal{S}_{S_n} in the case of F_n^*

Very similar! But:

Main difference: Definition of $\text{Hom}(\pi, \sigma)$.

For this: consider S_n inside the **partition monoid**.

Define $\text{Hom}(\pi, \sigma)$ as the set of all partitions which **contain** both π and σ .

Do the same as above to construct a fiat 2-category \mathcal{D} .

Theorem. Grothendieck decategorification of \mathcal{D} is isomorphic to $\mathbb{C}[F_n^*]$.

THANK YOU!!!